

Quantum Computers and Decoherence: Exorcising the Demon from the Machine

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ABSTRACT

Decoherence is the main obstacle to the realization of quantum computers. Until recently it was thought that quantum error correcting codes are the only complete solution to the decoherence problem. Here we present an alternative that is based on a combination of a decoherence-free subspace encoding and the application of strong and fast pulses: “encoded recoupling and decoupling” (ERD). This alternative has the advantage of lower encoding overhead (as few as two physical qubits per logical qubit suffice), and direct application to a number of promising proposals for the experimental realization of quantum computers.

Keywords: Quantum computers, decoherence, decoherence-free subspaces, dynamical decoupling, quantum error correction, selective recoupling, encoded recoupling and decoupling

1. INTRODUCTION

In the quest to construct large-scale quantum information processors, in particular quantum computers, decoherence is still the main obstacle to realization. Decoherence is the degradation of quantum information due to inevitable interactions with the environment. Early skepticism^{1,2} concerning the viability of quantum computation (QC) in the presence of decoherence was overcome by the discovery of quantum error correcting codes (QECCs),³⁻⁹ that flourished into a comprehensive theory that incorporates all elements of quantum computation.¹⁰ A QECC relies on an encoding of quantum information into the state of several quantum bits (qubits), and a closed-loop, active error diagnosis and correction procedure.¹¹ In principle, it is possible to correct arbitrary errors using sufficiently large QECCs.⁸ In practice, however, this may require a very large overhead, especially in terms of qubit resources.⁹ This is troubling in light of the substantial difficulties associated with generating and controlling systems with very large numbers of qubits. In addition, the theory of QECCs is rather abstract, in that it presumes that one can execute certain logical operations, but does not refer to the underlying Hamiltonians governing specific physical systems. Hence it is of interest to explore alternatives to QECCs, that are more economical in qubit resources, and that are directly tailored to specific quantum computer proposals.

Here we review our recent progress in developing such an alternative, “encoded recoupling and decoupling” (ERD),¹²⁻²¹ and report on some new results. ERD is based on a combination of encoding quantum information into a decoherence-free subspace (DFS)²²⁻²⁵ (for a review see Ref. 26), and the application of fast and strong dynamical-decoupling (or “bang-bang”, BB) pulses.^{15, 18, 19, 21, 27-45} The DFS encoding provides a first layer of protection against decoherence, while the BB pulses are used to efficiently reduce the remaining decoherence. The encoded “recoupling” part of ERD refers to the application of control operations that enact universal quantum logic, in a manner that is fully compatible with the DFS encoding and the BB pulses, and takes into account the experimentally available control resources, such as the underlying system Hamiltonian.¹³ The utility of ERD as a general method for quantum simulation, universal QC, and decoherence suppression has also been stressed and explored by Viola.³⁷ Our work builds in part on earlier efforts to combine universal QC with DFS encoding (without BB pulses).^{32, 35, 46, 47}

The structure of the paper is as follows. We begin in Section 2 with a brief formal summary of decoherence. We then give in Section 3 a simple example of a DFS encoding, protecting against collective dephasing. Section 4 shows how to perform universal QC on this encoding. We then briefly and formally review the dynamical decoupling method, in Section 5. The methods are then combined in Sections 6,7, where we show how BB pulses can eliminate decoherence sources beyond collective dephasing. In Section 8 we present a new and somewhat surprising result: the elimination of off-resonant transitions, that will typically be induced by BB pulses, via BB

pulses. Section 9 shows how to combine all the pieces together, by including logic gates with the BB pulses. Concluding remarks are presented in section 10.

2. DECOHERENCE

The dynamics of an open quantum system coupled to a bath is formally obtained from the time-ordered evolution

$$U(t) = \mathcal{T} \exp(-i \int^t H(t') dt') \quad (1)$$

under the combined system-bath Hamiltonian

$$\begin{aligned} H &= (H_S + H_C) \otimes I_B + I_S \otimes H_B + H_{SB} \\ H_{SB} &= \sum_{\gamma} S_{\gamma} \otimes B_{\gamma}, \end{aligned} \quad (2)$$

where I is the identity operator, H_S (H_B) is the internal Hamiltonian for the system (bath) alone, H_C is an externally applied control Hamiltonian, H_{SB} is the system-bath interaction Hamiltonian, and the S_{γ} (B_{γ}) are operators acting on the system (bath). The S_{γ} play the role of error generators in QC. Making the standard assumption of initially decoupled system and bath, one traces over the bath degrees of freedom in order to obtain the time-evolved reduced system density matrix:

$$\rho(t) = \text{Tr}_B[U(t) (\rho(0) \otimes \rho_B(0)) U^{\dagger}(t)], \quad (3)$$

where $\rho(0)$ [$\rho_B(0)$] is the initial density matrix of the system [bath]. As is well-known, $\rho(t)$ is in general a mixed state: $\text{Tr}[\rho(t)^2] < 1$, indicating that the environment has decohered the system (see, e.g., Ref. 48).

3. SIMPLE EXAMPLE OF DECOHERENCE-FREE SUBSPACES

We briefly review the simplest example of a DFS. Suppose that a system of K qubits (two-level systems) is coupled to a bath in a symmetric way, and undergoes a dephasing process. This can be described by the system-bath Hamiltonian

$$H_{SB}^{\text{coll-deph}} = \sum_{j=1}^K \sigma_j^z \otimes B^z \quad (4)$$

where σ_j^z is the Pauli- z matrix acting on the j th qubit and B^z is an arbitrary bath operator. Under this interaction qubit j undergoes the transformation

$$|0\rangle_j \rightarrow |0\rangle_j \quad |1\rangle_j \rightarrow e^{i\phi} |1\rangle_j, \quad (5)$$

which (after Tr_B) puts a random phase ϕ between the basis states $|0\rangle$ and $|1\rangle$ (eigenstates of σ_z with respective eigenvalues $+1$ and -1). This can also be described by the matrix $R_z(\phi) = \text{diag}(1, e^{i\phi})$ acting on the $\{|0\rangle, |1\rangle\}$ basis. As implied by (4), we assume that the phase has no space (j) dependence, i.e., the dephasing process is invariant under qubit permutations. Since the errors can be expressed in terms of the single Pauli spin matrix σ_z of the two-level system, this example is referred to as “collective phase damping”, or “weak collective decoherence”.⁴⁷ The more general situation when errors involving all three Pauli matrices are present, i.e., dissipation and dephasing, is referred to as “strong collective decoherence”,⁴⁷ or just “collective decoherence”.^{22, 24, 26} Without encoding a qubit initially in an arbitrary pure state $|\psi\rangle_j = a|0\rangle_j + b|1\rangle_j$ will decohere. This can be seen by calculating its density matrix as an average over all possible values of ϕ ,

$$\rho_j = \int_{-\infty}^{\infty} R_z(\phi) |\psi\rangle_j \langle \psi| R_z^{\dagger}(\phi) p(\phi) d\phi, \quad (6)$$

where $p(\phi)$ is a probability density, and we assume the initial state of all qubits to be a product state. For a Gaussian distribution, $p(\phi) = (4\pi\alpha)^{-1/2} \exp(-\phi^2/4\alpha)$, it is simple to check that

$$\rho_j = \begin{pmatrix} |a|^2 & ab^* e^{-\alpha} \\ a^* b e^{-\alpha} & |b|^2 \end{pmatrix}. \quad (7)$$

The decay of the off-diagonal elements in the computational basis is a signature of decoherence.

Let us now consider what happens in the two-qubit Hilbert space. The four basis states undergo the transformation

$$|0\rangle_1 \otimes |0\rangle_2 \rightarrow |0\rangle_1 \otimes |0\rangle_2 \quad (8)$$

$$|0\rangle_1 \otimes |1\rangle_2 \rightarrow e^{i\phi} |0\rangle_1 \otimes |1\rangle_2 \quad (9)$$

$$|1\rangle_1 \otimes |0\rangle_2 \rightarrow e^{i\phi} |1\rangle_1 \otimes |0\rangle_2 \quad (10)$$

$$|1\rangle_1 \otimes |1\rangle_2 \rightarrow e^{2i\phi} |1\rangle_1 \otimes |1\rangle_2. \quad (11)$$

Observe that the basis states $|0\rangle_1 \otimes |1\rangle_2$ and $|1\rangle_1 \otimes |0\rangle_2$ acquire the same phase. This suggests that a simple encoding trick can solve the decoherence problem. Let us define encoded states by $|0_L\rangle = |0\rangle_1 \otimes |1\rangle_2 \equiv |01\rangle$ and $|1_L\rangle = |10\rangle$. Then the state $|\psi_L\rangle = a|0_L\rangle + b|1_L\rangle$ evolves under the dephasing process as

$$|\psi_L\rangle \rightarrow a|0\rangle_1 \otimes e^{i\phi}|1\rangle_2 + b e^{i\phi}|1\rangle_1 \otimes |0\rangle_2 = e^{i\phi} |\psi_L\rangle, \quad (12)$$

and the overall phase thus acquired is clearly unimportant. This means that the 2-dimensional subspace $DFS_2(0) = \text{Span}\{|01\rangle, |10\rangle\}$ of the 4-dimensional Hilbert space of two qubits is *decoherence-free* (DF). The subspaces $DFS_2(2) = \text{Span}\{|00\rangle\}$ and $DFS_2(-2) = \text{Span}\{|11\rangle\}$ are also (trivially) DF, since they each acquire a global phase as well, 1 and $e^{2i\phi}$ respectively. Since the phases acquired by the different subspaces differ, there is decoherence *between* the subspaces.

For $K = 3$ qubits a similar calculation reveals that the subspaces $DFS_3(1) = \text{Span}\{|001\rangle, |010\rangle, |100\rangle\}$ and $DFS_3(-1) = \text{Span}\{|011\rangle, |101\rangle, |110\rangle\}$ are DF, as well the (trivial) subspaces $DFS_3(3) = \text{Span}\{|000\rangle\}$ and $DFS_3(3) = \text{Span}\{|111\rangle\}$.

More generally, let

$$\lambda_K = \text{number of } 0\text{'s minus the number of } 1\text{'s} \quad (13)$$

in a computational basis state (i.e., a bitstring) over K qubits. Then it is easy to check that any subspace spanned by states with constant λ_K is DF, and can be denoted $DFS_K(\lambda_K)$ in accordance with the notation above. The dimensions of these subspaces are given by the binomial coefficients: $d \equiv \dim[DFS_K(\lambda_K)] = \binom{K}{\lambda_K}$ and they each encode $\log_2 d$ qubits.

The encoding for the ‘‘collective phase damping’’ model discussed here has been tested experimentally. The first-ever experimental implementation of DFSs used the $DFS_2(0)$ subspace to protect against artificially induced decoherence in a linear optics setting.⁴⁹ The same encoding was subsequently used to alleviate the problem of external fluctuating magnetic fields in an ion trap quantum computing experiment,⁵⁰ and figures prominently in theoretical constructions of encoded, universal QC.^{13, 16, 21}

4. ENCODED UNIVERSAL LOGIC GATES

Next we show how to perform universal QC on the encoding $DFS_2(0) = \text{Span}\{|01\rangle, |10\rangle\}$.

4.1. Single encoded-qubit gates

Let X_i, Y_i, Z_i denote the standard Pauli matrices $\sigma_i^x, \sigma_i^y, \sigma_i^z$, acting on the i th physical qubit (we will use both notations interchangeably). In Ref. 12 it was shown that for the code $\{|0_L\rangle = |01\rangle, |1_L\rangle = |10\rangle\}$ the encoded logical operations (involving the first two physical qubits) are

$$\begin{aligned} \bar{X}_{12} &= \frac{1}{2}(X_1 X_2 + Y_1 Y_2), \\ \bar{Y}_{12} &= \frac{1}{2}(Y_1 X_2 - X_1 Y_2), \\ \bar{Z}_{12} &= \frac{1}{2}(Z_1 - Z_2). \end{aligned} \quad (14)$$

These operations form an $su(2)$ algebra (i.e., we think of them as Hamiltonians rather than unitary operators). We use a bar to denote logical operations on the encoded qubits. In Refs.^{12–14} these logical operations were denoted by T_{12}^α , $\alpha \in \{x, y, z\}$, and a detailed analysis was given on how to use typical solid-state Hamiltonians (Heisenberg, XXZ , and XY models, defined below) to implement quantum logic operations using this DFS encoding. E.g., the term $X_1X_2 + Y_1Y_2$ is the spin-spin interaction in the XY model, and $Z_1 - Z_2$ represents a Zeeman splitting. A static Zeeman splitting and a controllable XY interaction can be used to generate a universal set of logic gates. Similar conclusions hold when the XY interaction is replaced by a Heisenberg^{13, 51, 52} or XXZ interaction,¹⁴ or even by a Heisenberg interaction that includes an anisotropic spin-orbit term.¹⁷ We remark that, as first shown in Refs. 46, 47, the various types of exchange interactions can be made universal also without any single-qubit terms (such as a Zeeman splitting), by encoding into three or more qubits,^{47, 53–56} a result that has been termed “encoded universality”.⁵⁷

The two-qubit gate can be expressed as follows:

$$\begin{aligned} U_{ij}(\theta, \phi_i, \phi_j) &\equiv \exp(i\theta X_{\phi_i} X_{\phi_j}) \\ &= \cos \theta I_i I_j + i \sin \theta X_{\phi_i} X_{\phi_j}, \end{aligned} \quad (15)$$

where

$$X_\phi \equiv X \cos \phi + Y \sin \phi, \quad (16)$$

and I is the identity operator. In the context of trapped-ion QC,⁵⁸ the phase ϕ_i is the phase of the driving laser at the i th qubit, while θ is proportional to the Rabi frequency, and can be set over a wide range of values.⁵⁹ Introducing the operators

$$\tilde{X}_{ij} \equiv \frac{1}{2}(X_i X_j - Y_i Y_j), \quad \tilde{Y}_{ij} \equiv \frac{1}{2}(Y_i X_j + X_i Y_j) \quad (17)$$

(denoted R_{ij}^x , R_{ij}^y respectively in Refs. 12–14) we can express

$$U_{ij}(\theta, \phi_i, \phi_j) = \cos \theta \bar{I} + i \sin \theta (\cos \Delta \phi_{ij} \bar{X}_{ij} + \sin \Delta \phi_{ij} \bar{Y}_{ij} + \cos \Phi_{ij} \tilde{X}_{ij} + \sin \Phi_{ij} \tilde{Y}_{ij}), \quad (18)$$

where $\Phi_{ij} = \phi_i + \phi_j$. It is simple to check that \tilde{X}_{ij} and \tilde{Y}_{ij} annihilate the code subspace $\{|0_L\rangle = |01\rangle, |1_L\rangle = |10\rangle\}$ and have non-trivial action (as encoded X and Y) on the orthogonal subspace $\{|00\rangle, |11\rangle\}$. Therefore, upon restriction to the DFS we can write:

$$\begin{aligned} U_{ij}(\theta, \phi_i, \phi_j) &\xrightarrow{\text{DFS}} \bar{U}_{ij}(\theta, \Delta \phi_{ij}) \\ &= \exp(i\theta \bar{X}_{\Delta \phi_{ij}}) = \cos \theta \bar{I} + i \sin \theta \bar{X}_{\Delta \phi_{ij}}. \end{aligned} \quad (19)$$

The fact that \bar{U}_{ij} depends only on the relative phase $\Delta \phi_{ij}$ is crucial in the trapped-ion context: this quantity can be controlled by adjusting the angle between the driving laser and the interatomic axis, as well as by small adjustments of the trap voltages (which, in turn, control the trap oscillation frequency, and hence the qubit spacing), whereas it is much harder to control the absolute phase ϕ_i ,^{60, 61} and hence also Φ_{ij} . This is why the code subspace $\{|01\rangle, |10\rangle\}$ enjoys a preferred status over the subspace $\{|00\rangle, |11\rangle\}$.

The Hamiltonians in Eq. (14) generate the logic operations (19):

$$\begin{aligned} \exp(i\theta \bar{X}_{12}) &= U_{12}(\theta, \phi, \phi) = \bar{U}_{12}(\theta, 0) \\ \exp(i\theta \bar{Y}_{12}) &= U_{12}(\theta, \phi, \phi + \frac{\pi}{2}) = \bar{U}_{12}(\theta, \frac{\pi}{2}) \\ \exp(i\theta \bar{Z}_{12}) &= \exp(i\frac{\pi}{4} \bar{Y}_{12}) \exp(i\theta \bar{X}_{12}) \exp(-i\frac{\pi}{4} \bar{Y}_{12}) \\ &= \bar{U}_{12}(\frac{\pi}{4}, \pi/2) \bar{U}_{12}(\theta, 0) \bar{U}_{12}(-\frac{\pi}{4}, \pi/2). \end{aligned} \quad (20)$$

The third line follows from the elementary operator identity

$$X_\phi = X \cos \phi + Y \sin \phi = e^{-i\phi Z/2} X e^{i\phi Z/2} \quad (21)$$

which holds for any $su(2)$ angular momentum set $\{X, Y, Z\}$, i.e., operators that satisfy the commutation relation $[X, Y] = 2iZ$ (and cyclic permutations thereof), in particular also the encoded operators $\{\bar{X}, \bar{Y}, \bar{Z}\}$.

4.2. Entangling gate between pairs of encoded qubits

It is a well-known requirement of QC that in order to enact a universal set of logic gates (that allow any unitary transformation to be implemented) it is sufficient to be able to implement all single-qubit operations (as in the previous subsection), and to entangle pairs of qubits.⁶² In this subsection we discuss the latter requirement.

In Ref. 61 the following unitary gate was introduced, which is particularly suitable for QC using trapped ions:

$$\begin{aligned}
 U_4 &= \exp(-i\frac{\pi}{4}X_{\phi_1}X_{\phi_2}X_{\phi_3}X_{\phi_4}) \\
 &= \frac{1}{\sqrt{2}}(I_1I_2I_3I_4 - iX_{\phi_1}X_{\phi_2}X_{\phi_3}X_{\phi_4}) \\
 &\stackrel{\text{DFS}}{\mapsto} \frac{1}{\sqrt{2}}(\bar{I}_{12}\bar{I}_{34} - i\bar{X}_{\Delta\phi_{12}}\bar{X}_{\Delta\phi_{34}}) \\
 &= \exp(-i\frac{\pi}{4}\bar{X}_{\Delta\phi_{12}}\bar{X}_{\Delta\phi_{34}}). \tag{22}
 \end{aligned}$$

This gate can be used to entangle two DFS-qubits. It involves simultaneous control over two phase differences $\Delta\phi_{12}, \Delta\phi_{34}$, and thus control over the motion of two pairs of qubits. The case $\Delta\phi_{12} = \Delta\phi_{34} = 0$ was used in Ref. 60 to demonstrate entanglement of four trapped-qubit qubits, but this choice is not unique.

For general exchange Hamiltonians a different method is required. We review the encoded recoupling method introduced in Ref. 13. The exchange interaction quite generally has the form

$$H_{\text{ex}} = \sum_{\alpha=x,y,z} \sum_{i<j} J_{ij}^{\alpha} \sigma_i^{\alpha} \sigma_j^{\alpha} \tag{23}$$

$$= \sum_{i<j} J_{ij}^{-} R_{ij}^x + J_{ij}^{+} T_{ij}^x + J_{ij}^z \sigma_i^z \sigma_j^z, \tag{24}$$

where

$$T_{ij}^x = \frac{1}{2}(\sigma_i^x \sigma_j^x + \sigma_i^y \sigma_j^y), \quad R_{ij}^x = \frac{1}{2}(\sigma_i^x \sigma_j^x - \sigma_i^y \sigma_j^y), \tag{25}$$

and $J_{ij}^{\pm} = J_{ij}^x \pm J_{ij}^y$.

The isotropic (Heisenberg) case corresponds to $J_{ij}^{\alpha} \equiv J_{ij}$. The XY model is the case $J_{ij}^x = J_{ij}^y, J_{ij}^z = 0$. The XXZ model is the case $J_{ij}^x = \pm J_{ij}^y \neq J_{ij}^z$. A summary of QC proposals that fall into each category can be found in Ref. 13. The free Hamiltonian is

$$H_0 = \sum_i \frac{1}{2} \varepsilon_i \sigma_i^z, \tag{26}$$

where ε_i is the single-particle spectrum. In each instance of H_{ex} (Heisenberg, XY, XXZ) one typically has control over only one type of parameter out of the set $\{J_{ij}^{\alpha}, \varepsilon_i\}$.

Let $J_m^{\alpha} \equiv J_{2m-1, 2m}^{\alpha}$ ($\alpha = z, \pm$), and $\epsilon_m^{\pm} \equiv (\varepsilon_{2m-1} \pm \varepsilon_{2m})/2$. Let A and B be two angular momentum operators satisfying $su(2)$ commutation relations. Then it follows from Eq. (21) that the operation of ‘‘conjugating by A ’’,

$$\begin{aligned}
 C_A \circ \exp(iB) &\equiv \exp(-iA\pi/2) \exp(iB) \exp(iA\pi/2) \\
 &= \exp(-iB) \tag{27}
 \end{aligned}$$

causes B 's sign to be flipped.

Consider now an XXZ-type Hamiltonian where the J_m^+ parameters are controllable but ϵ_m^- and J_m^z are fixed. For simplicity let us consider just the case $J_m^- = 0$. Then we can rewrite

$$H = H_0 + H_{\text{ex}} = \sum_{m=1}^{N/2} \epsilon_m^- T_m^z - J_m^z T_m^z T_{m+1}^z + J_m^+ T_m^x, \tag{28}$$

where we have omitted a constant term. The important point is now that T_m^x and T_m^z satisfy $su(2)$ commutation relations. Therefore by using recoupling through “conjugation by T_m^x ” we can selectively turn on and off the single-encoded-qubit rotation T_m^z and the encoded-Ising interaction $T_m^z T_{m+1}^z$. This example of “encoded selective recoupling” establishes that encoded universal computation in the XXZ model can be done using control over the J_m^+ parameters alone.

Next, consider the XY model, with controllable J_{ij}^+ , but fixed ε_i . To implement encoded single-qubit operations, we can use the same encoded recoupling method as for the XXZ model. As for encoded two-qubit operations, we now no longer have the $\sigma_i^z \sigma_j^z$ terms. Since the XY model with nearest-neighbor interactions can be shown not to be universal,¹⁴ we turn on also next-nearest neighbor J_{ij}^+ terms (these can still be nearest-neighbor in a 2D hexagonal geometry). First note that $C_{T_{12}^x} \circ T_{23}^x = i\sigma_1^z \sigma_2^z T_{13}^x$. Now assume we can control J_{13}^+ ; then, using conjugation by $\pi/4$: $C_{\frac{1}{2}T_{13}^x} \circ (C_{T_{12}^x} \circ T_{23}^x) = \sigma_2^z (\sigma_3^z - \sigma_1^z)/2$. Since $\sigma_1^z \sigma_2^z$ is constant on the code subspace it can be ignored. On the other hand, $\sigma_2^z \sigma_3^z$ again acts as $-T_1^z T_2^z$, i.e., as an encoded $\sigma^z \otimes \sigma^z$. This establishes universal encoded computation in the XY model.

Taken together, the results in this section show how universal QC can be implemented using qubits encoded into a DFS offering protection against collective dephasing, while using only reasonable models of physically controllable Hamiltonians. We now move on to a discussion of how to reduce additional sources of decoherence.

5. DYNAMICAL DECOUPLING PULSES

We briefly review the dynamical decoupling method, introduced in Ref. 27, and further developed in Refs. 15, 18, 19, 21, 28–45. A set of “symmetrization” operations is chosen such that they form a discrete subgroup of the full unitary group of operations on the Hilbert space of the system. Denote this group $\mathcal{G} = \{V_j\}_{j=1}^{|\mathcal{G}|}$ where $V_1 = I$. The cycle time is $T_c = |\mathcal{G}|\Delta t$, where $|\mathcal{G}|$ is the number of symmetrization operations, and Δt is the time that the system evolves freely between operations under $U_0(t) \equiv \exp(-iH_S t)$. The symmetrized evolution after a single cycle is given by

$$U(T_c) = \prod_{j=1}^{|\mathcal{G}|} V_j^\dagger U_0(\Delta t) V_j \equiv e^{iH_{\text{eff}} T_c},$$

where the evolution under $H_{SB} + H_B$ has been neglected during application of the pulses V_j , which quantifies the sense in which these pulses have to be fast and strong (for details see, e.g, Refs. 21, 27, 30, 44, 45). H_{eff} denotes the resulting effective Hamiltonian, that can be computed to any order using the Magnus expansion.⁶³ In the BB limit one is interested in the evolution $U(NT_c)$ after N cycles, such that $N \rightarrow \infty$ and $\Delta t \rightarrow 0$ while $N\Delta t$ is finite. As shown in Ref. 34 one can then retain only the first order Magnus term, approximating H_{eff} by

$$H_{\text{eff}} = \frac{1}{|\mathcal{G}|} \sum_{j=1}^{|\mathcal{G}|} V_j^\dagger H V_j \equiv \Pi_{\mathcal{G}}(H). \quad (29)$$

The map $\Pi_{\mathcal{G}}$ is the projector into the centralizer Z of \mathcal{G} , defined as

$$Z(\mathcal{G}) = \{X \mid V_j^\dagger X V_j = X, \forall V_j \in \mathcal{G}\}.$$

Since the pulses g_j are unitary the centralizer equals the commutant in this case:

$$Z(\mathcal{G}) = \{X \mid [X, V_j] = 0, \forall V_j \in \mathcal{G}\}.$$

It is clear that $\Pi_{\mathcal{G}}$ commutes with the adjoint action $V_j \cdot V_j^\dagger$ for all j , so that if \mathcal{G} is generated by $\{I, H_S, S_\gamma\}$, the evolution will proceed without the operators S_γ affecting the system as errors: the evolution has been *symmetrized* with respect to \mathcal{G} .³¹ For a geometric interpretation of this symmetrization see Ref. 41, where it was also pointed out that while the BB pulses should have the effect of rotating H to the vertices of a symmetric object, they need in fact not form a group.

6. CREATING COLLECTIVE DEPHASING CONDITIONS USING DECOUPLING PULSES: REDUCING DECOHERENCE DURING STORAGE

One of the important advantages of the DFS encoding $\{|01\rangle, |10\rangle\}$ is that it is immune to collective dephasing. However, other sources of decoherence inevitably remain. In this and the following section, we algebraically classify all additional decoherence effects and show how they can be eliminated.

6.1. Creating collective dephasing on a pair of qubits

First, let us analyze the effect of breaking the collective dephasing symmetry, by considering a system-bath interaction of the form

$$H_{SB}^{\text{deph}(2)} = Z_1 \otimes B_1^z + Z_2 \otimes B_2^z \quad (30)$$

where B_1^z, B_2^z are arbitrary bath operators [compare to Eq. (4)]. This describes a general dephasing interaction on two qubits. The source of such dephasing during storage can be long wavelength, randomly fluctuating ambient magnetic fields,⁵⁰ that randomly shift the relative phase between the qubit $|0\rangle$ and $|1\rangle$ states through the Zeeman effect. The interaction can be rewritten as a sum over a collective dephasing term $Z_1 + Z_2$ and another, differential dephasing term $Z_1 - Z_2$, that is responsible for errors on the DFS:

$$H_{SB}^{\text{deph}(2)} = (Z_1 + Z_2) \otimes B_{\text{col}}^z + (Z_1 - Z_2) \otimes B_{\text{dif}}^z. \quad (31)$$

Here $B_{\text{col}}^z = (B_1^z + B_2^z)/2$ and $B_{\text{dif}}^z = (B_1^z - B_2^z)/2$. If B_{dif}^z were zero then there would only be collective dephasing and the DFS encoding would offer perfect protection. However, in general $B_{\text{dif}}^z \neq 0$, and the DFS encoding will not suffice to offer complete protection.

The crucial observation is that, since $Z_1 - Z_2 \propto \bar{X}_{12}$ [recall Eq. (14)], the offending term causes *logical* errors on the DFS.¹⁶ As shown in Refs. 15,37, then the problem of $B_{\text{dif}}^z \neq 0$ can be solved using a series of pulses that symmetrize $H_{SB}^{\text{deph}(2)}$ such that only the collective term remains. To do so note that since the offending term $\propto \bar{X}_{12}$, it anticommutes with $\bar{X}_{12} = \frac{1}{2}(X_1 X_2 + Y_1 Y_2)$. At the same time \bar{X}_{12} commutes with $Z_1 + Z_2$. This allows us to flip the sign of the offending term by using a pair of $\pm\pi/2$ pulses in \bar{X}_{12} , while leaving only the collective term. Evolution with the flipped sign followed by unaltered evolution leads to cancellation of the offending term. Specifically:¹⁵

$$e^{-iH_{SB}\tau} e^{-i\frac{\pi}{2}\bar{X}_{12}} e^{-iH_{SB}\tau} e^{i\frac{\pi}{2}\bar{X}_{12}} = e^{-i(Z_1+Z_2)\otimes B_{\text{col}}^z 2\tau}, \quad (32)$$

or, in terms of gates:

$$e^{-iH_{SB}\tau} \bar{U}_{12}\left(-\frac{\pi}{2}, 0\right) e^{-iH_{SB}\tau} \bar{U}_{12}\left(\frac{\pi}{2}, 0\right) = e^{-i(Z_1+Z_2)\otimes B_{\text{col}}^z 2\tau}, \quad (33)$$

where $\bar{U}_{ij}(\theta, \Delta\phi_{ij})$ was defined in Eq. (19). This equation means that the system-bath coupling effectively looks like collective dephasing at the end of the pulse sequence. Thus, the system is periodically (every 2τ) projected into the DFS.

In order for the the procedure described in Eq. (33) to work, the gate $\bar{U}_{12}(\pm\frac{\pi}{2}, 0)$ must be executed at a timescale faster than the cutoff frequency associated with the fluctuating (magnetic) fields causing the differential dephasing term in $H_{SB}^{\text{deph}(2)}$.

6.2. Creating collective dephasing on a block of four qubits

So far we have discussed creation of collective dephasing conditions on a single DFS qubit. However, it is essential for the reliable execution of an entangling logic gate to have collective dephasing over all four qubits participating in the gate, even if only two are coupled at a time. A procedure for creating collective *decoherence* conditions over blocks of 3, 4, 6 and 8 qubits was given in Ref. 15. Here we show how to do the same for a block of 4 qubits with collective dephasing.

Let us start with a general dephasing Hamiltonian on N qubits, and rewrite it in terms of nearest-neighbor sums and differences:

$$H_{SB}^{\text{deph}} = \sum_{i=1}^N Z_i \otimes B_i \quad (34)$$

$$= \sum_{j=1}^{N/2} (Z_{2j} + Z_{2j-1}) \otimes B_{2j}^+ + (Z_{2j} - Z_{2j-1}) \otimes B_{2j}^-, \quad (35)$$

where $B_{2j}^{\pm} \equiv (B_{2j} \pm B_{2j-1})/2$. As noted above, $Z_{2j} - Z_{2j-1} \propto \bar{X}_{2j-1,2j}$, so that to eliminate all nearest-neighbor differences of the form $(Z_{2j} - Z_{2j-1})$ we can use the collective decoupling pulse $X_{nn} = \bigotimes_{j=1}^{N/2} e^{i\frac{\pi}{2}\bar{X}_{2j-1,2j}}$:

$$e^{-iH_{SB}\tau} X_{nn} e^{-iH_{SB}\tau} X_{nn}^\dagger = e^{-i2\tau \sum_{j=1}^{N/2} (Z_{2j} + Z_{2j-1}) \otimes B_{2j}^+}, \quad (36)$$

or, in gate terms:

$$e^{-iH_{SB}\tau} \left[\bigotimes_{j=1}^{N/2} \bar{U}_{2j-1,2j}(-\frac{\pi}{2}, 0) \right] e^{-iH_{SB}\tau} \left[\bigotimes_{j=1}^{N/2} \bar{U}_{2j-1,2j}(\frac{\pi}{2}, 0) \right] = e^{-i2\tau \sum_{j=1}^{N/2} (Z_{2j} + Z_{2j-1}) \otimes B_{2j}^+}.$$

The next step is to eliminate next-nearest neighbor differential terms. To this end let us rewrite the outcome of the X_{nn} pulse in terms of sums and differences over blocks of four ions:

$$\begin{aligned} \sum_{j=1}^{N/2} (Z_{2j} + Z_{2j-1}) \otimes B_{2j}^+ &= \sum_{j=1}^{N/2} [Z_{2j+2} + Z_{2j+1} + Z_{2j} + Z_{2j-1}] \otimes B_{2j}^{+,+} \\ &+ \sum_{j=1}^{N/2} [(Z_{2j+2} - Z_{2j}) + (Z_{2j+1} - Z_{2j-1})] \otimes B_{2j}^{+,-}, \end{aligned} \quad (37)$$

where $B_{2j}^{+,\pm} \equiv (B_{2j+2}^+ \pm B_{2j}^+)/2$. The term in the first line contains only the desired block-collective dephasing over 4 ions. The term in the second line contains undesired differential dephasing terms that we wish to eliminate. But these terms once again have the appearance of encoded Z operators, between next-nearest neighbor ion pairs. Therefore we need to apply a second collective pulse $X_{nnn} = \bigotimes_{j=1}^{N/2} e^{i\frac{\pi}{2}\bar{X}_{2j-1,2j+1}} e^{i\frac{\pi}{2}\bar{X}_{2j,2j+2}}$, that applies encoded X operators on these qubit pairs. At this point we are left just with collective dephasing terms on blocks of 4 qubits, as required:

$$\begin{aligned} e^{-i2\tau \sum_{j=1}^{N/2} (Z_{2j} + Z_{2j-1}) \otimes B_{2j}^+} \left[\bigotimes_{j=1}^{N/2} \bar{U}_{2j-1,2j+1}(-\frac{\pi}{2}, 0) \bar{U}_{2j,2j+2}(-\frac{\pi}{2}, 0) \right] &\times \\ e^{-i2\tau \sum_{j=1}^{N/2} (Z_{2j} + Z_{2j-1}) \otimes B_{2j}^+} \left[\bigotimes_{j=1}^{N/2} \bar{U}_{2j-1,2j+1}(\frac{\pi}{2}, 0) \bar{U}_{2j,2j+2}(\frac{\pi}{2}, 0) \right] &= e^{-i4\tau \sum_{j=1}^{N/2} (Z_{2j+2} + Z_{2j+1} + Z_{2j} + Z_{2j-1}) \otimes B_{2j}^{+,+}}. \end{aligned} \quad (38)$$

This pulse sequence is important to ensure that collective dephasing conditions will prevail during the execution of logic gates between DFS qubits.

7. REDUCTION OF ALL REMAINING DECOHERENCE ON A SINGLE DFS QUBIT DURING LOGIC GATE EXECUTION

The reduction of differential dephasing errors, as in the previous subsection, is particularly relevant for storage errors. However, this is only the first step. Additional sources of decoherence may take place during storage,

and in particular during the execution of logic gates. It is useful to provide a complete algebraic classification of the possible decoherence processes. This will allow us to see what can be done using decoupling pulses. To this end let us now write the system-bath Hamiltonian on two physical qubits in the general form

$$H_{SB} = H_{\text{Leak}} + H_{\text{Logi}} + H_{\text{DFS}} \quad (39)$$

where

$$\begin{aligned} H_{\text{DFS}} &= \text{Span}\left\{\frac{ZI + IZ}{2}, \frac{XY + YX}{2}, \frac{XX - YY}{2}, ZZ, II\right\} \\ H_{\text{Leak}} &= \text{Span}\{XI, IX, YI, IY, XZ, ZX, YZ, ZY\} \\ H_{\text{Logi}} &= \text{Span}\left\{\bar{X} = \frac{XX + YY}{2}, \bar{Y} = \frac{YX - XY}{2}, \bar{Z} = \frac{ZI - IZ}{2}\right\} \end{aligned} \quad (40)$$

where I is the identity operator, $XZ \equiv X_1 Z_2$ (etc.), and where Span means a linear combination of these operators tensored with bath operators. The 16 operators in Eq. (40) form a complete basis for all 2-qubit operators. This classification, first introduced in Ref. 16, has the following significance. The operators in H_{DFS} either vanish on the DFS, or are proportional to identity on it. In either case their effect is to generate an overall phase on the DFS, so they can be safely ignored from now on. The operators in H_{Leak} are the *leakage errors*: terms that cause transitions between states inside and outside of the DFS. A universal and efficient decoupling method for eliminating such errors, for arbitrary numbers of (encoded) qubits was given in Ref. 18. Finally, the operators in H_{Logi} have the form of logic gates on the DFS. However, these are undesired logic operations, since they are coupled to the bath, and thus cause decoherence.

In the previous subsection we showed how to eliminate the logical error \bar{Z} , but we see now that this was only one error in a much larger set. To deal with the additional errors it is useful at this point to introduce a more compact notation for the pulse sequences. We denote by $[\tau]$ a period of evolution under the free Hamiltonian, i.e., $U(\tau) \equiv \exp(-iH_{SB}\tau) \equiv [\tau]$, and further denote

$$P \equiv \bar{U}_{12}\left(-\frac{\pi}{2}, 0\right) = \exp\left(-i\frac{\pi}{2}\bar{X}_{12}\right). \quad (41)$$

Thus Eq. (33) can be written as:

$$\exp[-i(B_1^z + B_2^z)(Z_1 + Z_2)\tau] = [\tau, P, \tau, P^\dagger]. \quad (42)$$

As a first step in dealing with the additional errors, note that the symmetrization procedure $[\tau, P, \tau, P^\dagger]$ can in fact achieve more than just the elimination of the differential dephasing $Z_1 - Z_2$ term. Since \bar{X}_{12} also anticommutes with $\bar{Y}_{12} = \frac{1}{2}(Y_1 X_2 - X_1 Y_2) \in H_{\text{Logi}}$, if such a term appears in the system-bath interaction it too will be eliminated using the same procedure.

So far we have used a $\frac{\pi}{2}\bar{X}_{12}$ pulse. Interestingly, the Hamiltonian \bar{X}_{12} can also be used to eliminate all leakage errors.¹⁶ To see this, note that $\bar{U}_{12}(\pm\pi, 0) = \exp(\pm i\pi\bar{X}_{12}) = Z_1 Z_2$. This operator anticommutes with *all* terms in H_{Leak} . Hence it too can be used in a parity-kick pulse sequence, that will eliminate all the leakage errors.

At this point we are left with just a single error: $\bar{X}_{12} \otimes B$ itself, in H_{Logi} . Clearly, we cannot use a pulse generated by \bar{X}_{12} to eliminate this error. Instead, to deal with this error we need to introduce one more pulse pair that anticommutes with \bar{X}_{12} , e.g., $\exp(\pm i\frac{\pi}{2}\bar{Y}_{12}) = \bar{U}_{12}(\pm\frac{\pi}{2}, \frac{\pi}{2})$.

Let us now see how to combine all the decoherence elimination pulses into one efficient sequence. First we introduce the abbreviations

$$\begin{aligned} \Pi &\equiv \bar{U}_{12}(\pm\pi, 0) = \exp(\pm i\pi\bar{X}_{12}) = \Pi^\dagger = PP \\ Q &\equiv \bar{U}_{12}\left(-\frac{\pi}{2}, \frac{\pi}{2}\right) = \exp\left(-i\frac{\pi}{2}\bar{Y}_{12}\right) \\ \Lambda &\equiv \bar{U}_{12}\left(\pm\pi, \frac{\pi}{2}\right) = \exp(\pm i\pi\bar{Y}_{12}) = \Lambda^\dagger = QQ \end{aligned} \quad (43)$$

As argued above, the π pulse Π eliminates H_{Leak} :

$$\exp[-i(H_{\text{Logi}} + H_{\text{DFS}})2\tau] = [\tau, \Pi, \tau, \Pi]. \quad (44)$$

Now let us discuss adding the extra pulses needed to achieve full decoherence elimination. The $\pi/2$ pulse P eliminates \bar{Y} and \bar{Z} in H_{Logi} . Combining this with the sequence for leakage elimination we have the sequence of 4 pulses:

$$\begin{aligned} e^{-i(H_{\text{DFS}} + \bar{X} \otimes B_{\bar{X}})4\tau} &= [U(\tau)\Pi U(\tau)\Pi]P^\dagger [U(\tau)\Pi U(\tau)\Pi]P \\ &= [\tau, \Pi, \tau, P, \tau, \Pi, \tau, P^\dagger], \end{aligned} \quad (45)$$

(where we have used $\Pi P^\dagger = P$, $\Pi P = P^\dagger$).

If we wish to entirely eliminate decoherence then we are left just with getting rid of the logical error due to \bar{X} . To eliminate it we now combine with the \bar{Y} -direction, $\pi/2$ pulse, Q :

$$\begin{aligned} e^{-iH_{\text{DFS}}8\tau} &= [U(\tau)\Pi U(\tau)PU(\tau)\Pi U(\tau)P^\dagger]Q^\dagger [U(\tau)\Pi U(\tau)PU(\tau)\Pi U(\tau)P^\dagger]Q \\ &= [\tau, \Pi, \tau, P, \tau, \Pi, \tau, P^\dagger, Q^\dagger, \tau, \Pi, \tau, P, \tau, \Pi, \tau, P^\dagger, Q] \end{aligned} \quad (46)$$

which takes ten pulses. Unfortunately it is not possible to compress this further, since $P^\dagger Q = (i\bar{X})(-i\bar{Y}) = i\bar{Z}$ and $P^\dagger Q^\dagger = -i\bar{Z}$, neither of which cannot be generated directly (in one step) from the available gate $\bar{U}_{ij}(\theta, \Delta\phi_{ij}) = \cos\theta\bar{I} + i\sin\theta\bar{X}_{\Delta\phi_{ij}}$.

8. ELIMINATION OF OFF-RESONANT TRANSITIONS

One important caveat in our discussion so far is that, because we need very strong and fast pulses, our gate operation may become imperfect. Specifically, off-resonant coupling may become important. This can cause *unitary* leakage errors from the DFS. These can in turn be reduced using the methods in Refs. 64, 65. Here we present an alternative and new method, that uses BB pulses generated in terms of the system Hamiltonian.

Consider an N -level system Hamiltonian

$$H_0 = \sum_{i=1}^N E_i |i\rangle \langle i|$$

where the levels $i = 1, 2$ denote our qubit. If all parameters are fixed, the evolution $U_0(t) = \exp(-iH_0t)$ is always on. If we turn on an interaction H_I , then the total Hamiltonian is $H = H_0 + H_I$, which can be written as

$$H = \sum_{i=1}^N H_{ii} |i\rangle \langle i| + \sum_{i>j>2}^N (H_{ij} |i\rangle \langle j| + h.c.) + H_L$$

where the leakage term is

$$H_L = |1\rangle \sum_{i=3}^N H_{1i} \langle i| + |2\rangle \sum_{i=3}^N H_{2i} \langle i| + h.c.$$

Now note that

$$U_0^\dagger(t) H U_0(t) = \sum_{i=1}^N H_{ii} |i\rangle \langle i| + \sum_{i>j=1}^N (e^{-i(E_i - E_j)t} H_{ij} |i\rangle \langle j| + h.c.).$$

Using this we now show how to eliminate each of the leakage terms one by one. First, we eliminate $|1\rangle H_{13} \langle 3| + h.c.$ by the BB sequence

$$2H' = H + U_0^\dagger\left(\frac{\pi}{E_1 - E_3}\right) H U_0\left(\frac{\pi}{E_1 - E_3}\right)$$

The new Hamiltonian H' does not contain $|1\rangle H_{13} \langle 3| + h.c.$ Next, we eliminate $|2\rangle H'_{23} \langle 3| + h.c.$ by $U_0(\frac{\pi}{E_2 - E_3})$, where $H'_{23} = \exp(-i\frac{\pi(E_2 - E_3)}{E_1 - E_3})H_{23}$. We can clearly repeat this procedure so as to eliminate all leakage. *The crucial point is that since we have used the system Hamiltonian to generate BB pulses, off-resonant transitions do not take place: H_0 has no matrix elements between different levels.*

A concern is how we to obtain $U_0^\dagger(t) = \exp(+iH_0t)$. This problem is shared by other methods dealing with the same problem.⁶⁵ In principle it can be solved provided the level spacings are rationally related. Note in this context that usually we do not need to eliminate all H_{1i} , since $H_{1i} \gg H_{1i+1}$ and typically decrease exponentially according to time-independent perturbation theory. We have assumed that there is no degeneracy. If there is, then all degenerate transitions will be eliminated simultaneously, so the procedure is simplified.

9. COMBINING LOGIC GATES WITH DECOUPLING PULSES

So far we have discussed computation using the encoded recoupling method (Section 4), and encoded decoupling (Sections 6,7). We now put the two together in order to obtain the full ERD picture. At least two methods are available for combining quantum computing operations with the sequences of decoupling pulses we have presented above. For a general analysis of this issue see Ref. 35.

9.1. Fast + Strong Gates Method

The decoupling pulse sequences given in Sec. 7 “stroboscopically” create collective dephasing conditions at the conclusion of each cycle. As noted above, this is equivalent to a periodic projection into the DFS. This property allows for “stroboscopic” quantum computation at the corresponding projection times.³⁵ Here the computation pulses need to be synchronized with the decoupling pulses, and inserted at the end of each cycle. Because of the conditions on the validity of the BB method,²⁷ the amount of time available for implementation of a logic gate is no more than the bath correlation time $\tau_c = 2\pi/\omega_c$. (An exception to this rule is the case of $1/f$ noise.⁴⁴) Assuming the dominant decoherence contributions not accounted for by the DFS encoding to come from differential dephasing (setting the τ_c time-scale), and given that we already assumed that we can use pulses with interval $\Delta t \ll \tau_c$, it is consistent to assume that we can then also perform logic gates on the same time scale.

9.2. Fast + Weak Gates Method

There may be an advantage to using fast but weak pulses for the logic gates, while preserving the fast + strong property of the decoupling pulses. To see how to combine logic gates with decoupling in this case, let us denote by $H_S = X_{\phi_i} X_{\phi_j}$ the controllable system Hamiltonian that generates the entangling gate $U_{ij}(\theta, \phi_i, \phi_j)$ [recall Eq. (15)]. Suppose first that we turn on this logic-gate generating Hamiltonian in a manner that is neither very strong nor very fast, so that the system-bath interaction is not negligible while H_S is on (this obviously puts less severe demands on experimental implementation). Then the corresponding unitary operator describing the dynamics of system plus bath is:

$$\tilde{U}(t) = \exp[-it(H_S + H_{SB} + H_B)]. \quad (47)$$

Now, if we choose H_S so that it commutes with the decoupling pulses, then we can show that after decoupling

$$\tilde{U}(t) \mapsto \exp[-i2t(H_S + H_B)], \quad (48)$$

provided t is sufficiently small. Tracing out the bath then leaves a purely unitary, decoherence-free evolution on the system. To prove this, assume we have chosen t' and the decoupling Hamiltonian H'_S so that (i) $\exp(-it'H'_S)H_{SB}\exp(it'H'_S) = -H_{SB}$, and (ii) $[H'_S, H_S] = 0$. Then

$$\tilde{U}(t)e^{-it'H'_S}\tilde{U}(t)e^{-it'H'_S} = \tilde{U}(t)e^{-it[H_S + e^{-it'H'_S}H_{SB}e^{it'H'_S} + H_B]} \quad (49)$$

$$= e^{-it(H_S + H_{SB} + H_B)} e^{-it(H_S - H_{SB} + H_B)} \quad (50)$$

$$= e^{-\{2it(H_S + H_B) + t^2([H_{SB}, H_S] + [H_{SB}, H_B]) + O(t^3)\}}, \quad (51)$$

where we have used the Baker-Campbell-Hausdorff formula, $\exp(\alpha A)\exp(\alpha B) = \exp\{\alpha(A + B) + \frac{\alpha^2}{2}[A, B] + O(\alpha^3)\}$.

Let us now show how to efficiently combine logic operations and decoupling pulses. For simplicity consider only the case where we can neglect the \bar{X} error, i.e., our decoupling sequence is the 4-pulse one given in Eq. (45). Suppose we wish to implement a logical X operation, i.e., $\exp(-i\theta\bar{X}_{12})$. Recall that this involves turning on the Hamiltonian $H_S^X = \Omega_X X_\phi X_\phi \xrightarrow{\text{DFS}} \Omega_X \bar{X}_{12}$ between two physical qubits. Because the decoupling pulses $P = \exp(-i\frac{\pi}{2}\bar{X}_{12})$ and $\Pi = \exp(\pm i\pi\bar{X}_{12})$ are generated in terms of the same Hamiltonian, they commute with H_S^X while eliminating H_{SB} (except for the terms in H_{SB} that have trivial action on the DFS). Thus the conditions under which Eq. (48) were shown to hold are satisfied. This allows us to insert the logic gates into the four free evolution periods involved in the pulse sequence of Eq. (45). Thus, the full pulse sequence that combines creation of collective dephasing conditions with execution of the logic gate is:

$$e^{-it(\Omega_X \bar{X}_{12} + H_{\text{DFS}})} = \tilde{U}(t/4)\Pi\tilde{U}(t/4)P\tilde{U}(t/4)\Pi\tilde{U}(t/4)P^\dagger, \quad (52)$$

with $\tilde{U}(t) = \exp[-it(H_S^X + H_{SB} + H_B)]$, and which, using the DFS encoding, is equivalent to the desired $\exp(-i\theta\bar{X}_{12})$. This involves 8 control pulses, 4 of which are of the fast+strong type (those involving P and Π), and 4 of which must be fast, but need not be so strong that we can neglect H_{SB} .

If we wish to implement logical Y operation, i.e., $\exp(-i\theta\bar{Y}_{12})$, then we cannot now use P and Π , since they anticommute with \bar{Y}_{12} and will eliminate it. Instead we should use decoupling pulses generated in terms of \bar{Y}_{12} , which will also have the desired effect of eliminating H_{Leak} , as well as \bar{X} and \bar{Z} logical errors, while commuting with the \bar{Y} logic operations (and for this reason can of course not eliminate \bar{Y} errors). These are just the Q and Λ pulses defined in Eq. (43). In gate terms this implies turning on the Hamiltonian $H_S^Y = \Omega_Y X_\phi X_{\phi+\pi/2} \xrightarrow{\text{DFS}} \Omega_Y \bar{Y}_{12}$ between two physical qubits. Thus:

$$e^{-it(\Omega_Y \bar{Y}_{12} + H_{\text{DFS}})} = \tilde{U}(t/4)\Lambda\tilde{U}(t/4)Q\tilde{U}(t/4)\Lambda\tilde{U}(t/4)Q^\dagger, \quad (53)$$

with $\tilde{U}(t) = \exp[-it(H_S^Y + H_{SB} + H_B)]$, and which, using the DFS encoding, is equivalent to the desired $\exp(-i\theta\bar{Y}_{12})$.

Finally, to generate single DFS-qubit rotations about an arbitrary axis we can combine Eqs. (52),(53) according to the Euler angles construction.⁶² Given that Eqs. (52),(53) each take 8 pulses, the Euler angle method will generate an arbitrary DFS-qubit rotation in at most 24 pulses.

Concerning gates that entangle two DFS-qubits, the situation may be more involved, since now the next-nearest neighbor pulses in Eq. (38), that create the collective dephasing conditions on four qubits, do not all commute with the U_4 gate of Eq. (22). Therefore here we must resort to the strong + fast method of the previous subsection, i.e., we need to synchronize the U_4 pulses with the end of the decoupling pulse sequence. We do not here analyze the situation with respect to the exchange Hamiltonian implementation of conditional logic gates.

Taken together, the methods described in this section provide an explicit way to implement universal QC in a manner that offers protection against all sources of qubit decoherence, using a fast + strong (or fast + weak) versions of logic gates.

10. DISCUSSION AND CONCLUSIONS

We have proposed a method of encoded recoupling and decoupling (ERD) for performing decoherence-protected quantum computation. Our method combines an encoding into qubit-pair decoherence-free subspaces (each pair yielding one encoded qubit), with sequences of recoupling and decoupling pulses. The qubit encoding protects against collective dephasing processes, while the decoupling pulses symmetrize all other sources of decoherence into a collective dephasing interaction. The recoupling pulses are used to implement encoded quantum logic gates, either during or in between the decoupling pulses.

The dynamical decoupling method requires an exponential number of pulses if the most general form of decoherence is to be suppressed, that can couple arbitrary numbers of qubits to the environment (total decoherence²⁴). This exponential scaling is avoided here by focusing on decoherence elimination inside blocks of *finite* size (e.g., at most four qubits) where arbitrary decoherence is allowed. However, we have implicitly assumed that there are no decoherence processes coupling different blocks. This is a reasonable assumption for most quantum

computer implementation, provided the different blocks can be kept sufficiently far apart until they need to be brought together in order to execute inter-block logic gates. When this happens, ERD can still be efficiently applied on the temporarily larger block.

It may be questioned whether there is any advantage in using ERD compared to methods of active quantum error correcting codes (QECC). Both ERD and QECC are capable of dealing with arbitrary decoherence processes, and are fully compatible with universal quantum computation. There are two main advantages to ERD: First, we need only two physical qubits per logical qubit, compared to a redundancy of five to one, in order to handle all single-qubit errors in QECC.⁶ So far the most advanced experiments outside NMR, i.e., those involving trapped ions, have used up to four qubits,⁶⁰ so that this encoding economy is a distinct advantage for near-term experiments. Second, our method is directly compatible with Hamiltonians describing a variety of quantum computer proposals. On the other hand it is not clear how to directly use QECC given Hamiltonians describing specific systems. These are general features of ERD: economy of encoding redundancy and use of only the most easily controllable interactions. The disadvantage of ERD compared to QECC is that there does not exist, at this point, a result analogous to the threshold theorem of fault tolerant quantum error correction. This means that we cannot yet guarantee full scalability of ERD as a stand-alone method, because we do not yet know how to compensate for imperfect pulses. However, in principle it is always possible to concatenate ERD with QECC, as done, e.g., for DFS with QECC in Refs. 66, and then the standard fault tolerance results apply. We expect that the theory of composite pulses⁶⁷ will also play a key role in this further development of ERD.

We hope that the methods proposed here will inspire experimentalists to implement encoded recoupling and decoupling in the lab, thus demonstrating the possibility of fully decoherence-protected quantum computation.

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